

Introduction and Overview

(Read Cow book Chapter 1)

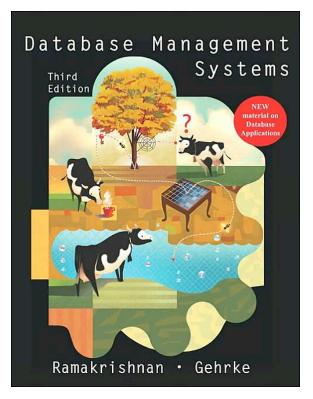
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Course Administrivia



- Book
 - Cow book
 - Somewhat Dense
 - Cover about 80%
- Instructor
 - Leonard McMillan
- Teaching Assistant
 - To be named
- When will we meet?
 - Mondays and Wednesdays (barring university holidays)





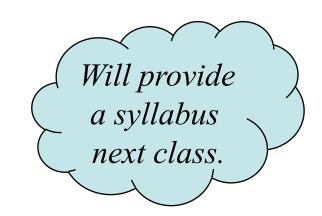


Course Logistics

Website (not up yet):

http://www.cs.unc.edu/Courses/comp521-f10

- look here first for
 - News, problem-set hints, and helpful resources
 - Revisions, solutions, and corrections to problem sets
- ✤ Office Hours: TBA
- Grading
 - 5 Problem sets (worth 9% each)
 - 2 Quizzes (worth 15% each) Final Exam (worth 25%)
- Problem Sets



Roughly one every two weeks, except weeks with quizzes





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35



- Relational Model
- Relational Algebra
- Relational Calculus
- Normal Forms



- Integrating Dbases & programs
- Web-based Dbase use
- Triggers and Active databases

ApplicationsSystems

Foundations

- Database Indexing
- Query Evaluations
- Query Optimization
- Transactions and Concurrency

Emphasis

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Where Databases fit into CS

Designing Programs

- Syntax
- Semantics
- Abstraction
- Designing Algorithms
 - Correctness
 - Efficiency
- Designing Data
 - Generalization
 - Portability
 - Independence

Data sets are growing far faster than either languages used to process them or the algorithms used to manage them.



- A very large, integrated collection of bits.
- Models real-world <u>enterprise</u>.
 - Entities (e.g., students, courses)
 - Relationships (e.g., Madonna is taking Comp 521)
- A <u>Database Management System (DBMS)</u> is a software package designed to store and manage databases.





- Application must stage large datasets between main memory and secondary storage (e.g., buffering, page-oriented access, 32-bit addressing, etc.)
- Special code for different queries
- Must protect data from inconsistency due to multiple concurrent users
- Crash recovery
- Security and access control



- * Data Independence
- Efficient access.
- Reduced application development time.
- Data integrity and security.
- Uniform data administration.
- Concurrent access, recovery from crashes.







- Shift from <u>computation</u> to <u>information</u>
 - at the "low end": dynamic web spaces
 - at the "high end": scientific applications
- Datasets increasing in diversity and volume.
 - Digital libraries, interactive video, Human Genome project, Earth-Observing Satellite (EOS) project
 - ... need for DBMS exploding
- DBMS encompasses most of CS
 - OS, languages, theory, AI, multimedia, logic







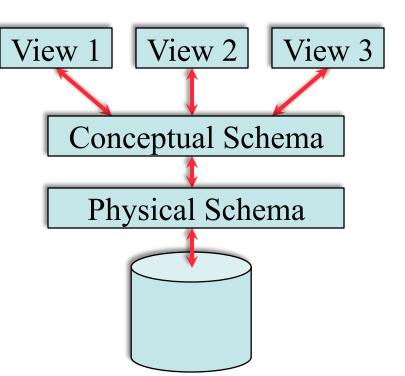
- A <u>data model</u> is a collection of concepts for describing data.
- A <u>schema</u> is a description of a particular collection of data, using the a given data model.
- The <u>relational model of data</u> is the most widely used model today.
 - Main concept: <u>relation</u>, basically a table with rows and columns.
 - Every relation has a <u>schema</u>, which describes the allowed contents of columns, or fields.





Levels of Abstraction

- Many <u>views</u>, single
 <u>conceptual (logical) schema</u>
 and <u>physical schema</u>.
 - Views describe how users see the data.
 - Conceptual schema defines logical structure
 - Physical schema describes the files and indexes used.



 Schemas are defined using a Data-Description Languages (DDLs); data is modified/queried using Data-Management Languages (DMLs).





Example: University Database

Conceptual schema:

- Students(sid: string, name: string, login: string, dob: date, gpa: real)
- Courses(cid: string, cname: string, credits: integer)
- Enrolled(sid: string, cid: string, grade: string)
- Physical schema:
 - Relations stored as unordered files.
 - Index on first column of Students.
- External Schema (View):
 - Course_info(cid: string, enrollment: integer)





- Data Independence*
- Applications insulated from how data is actually structured and stored.
- Logical data independence: Protection from changes in logical structure of data.
- ✤ <u>Physical data independence</u>: Protection from changes in *physical* structure of data.

► One of the most important benefits of using a DBMS!



- Concurrent execution of multiple user queries is essential for good DBMS performance.
 - Because disk accesses are frequent, and relatively slow, it is important to keep the cpu humming by working on several user programs concurrently.
- Interleaving actions of different user programs can lead to inconsistency: e.g., check is cleared while account balance is being computed.
- DBMS ensures such problems don't arise: users can pretend they are using a single-user system.



Database Transactions



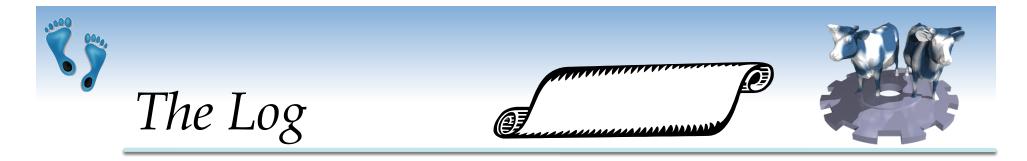
- Key concept is <u>transaction (Xact</u>), which is an <u>atomic</u> sequence of database actions.
- Each transaction, executed completely, must leave the DB in a <u>consistent state</u> if DB is consistent when the transaction begins.
 - Users can specify some simple <u>integrity constraints</u> on the data, and the DBMS will enforce these constraints.
 - Beyond this, the DBMS does not really understand the semantics of the data. (e.g., it does not understand how the interest on a bank account is computed).
 - Thus, ensuring that a transaction (run alone) preserves consistency is ultimately the user's responsibility!

Scheduling Concurrent Transactions

- DBMS ensures that execution of {T1, ..., Tn} is equivalent to some <u>serial</u> execution T1' ... Tn'.
 - Before reading/writing an object, a transaction requests a lock on the object, and waits till the DBMS gives it the lock. All locks are released at the end of the transaction. (Strict Two-Phase Locking (2PL) protocol.)
 - Idea: If an action of Ti (say, writing X) affects Tj (which perhaps reads X), one of them, say Ti, will obtain the lock on X first and Tj is forced to wait until Ti completes; this effectively orders the transactions.
 - What if Tj already has a lock on Y and Ti later requests a lock on Y? (<u>Deadlock</u>!) Ti or Tj is <u>aborted</u> and restarted!



- DBMS ensure *atomicity* (all-or-nothing property) even if system crashes in the middle of a Xact.
- Idea: Keep a <u>log</u> (history) of all actions carried out by the DBMS while executing a set of Xacts:
 - Before a change is made to the database, the corresponding log entry is forced to a safe location. (Write-Ahead Log (<u>WAL</u>) <u>protocol</u>)
 - After a crash, the effects of partially executed transactions are <u>undone</u> using the log. (Thanks to WAL, if log entry wasn't saved before the crash, corresponding change was not applied to database!)



- The following actions are recorded in the log:
 - *Ti writes an object*: The old value and the new value.
 - Log record must go to disk *before* the changed page!
 - *Ti commits/aborts*: A log record indicating this action.
- Log records chained together by Xact id, so it's easy to undo a specific Xact (e.g., to resolve a deadlock).
- Log is often *duplexed* and *archived* on "stable" storage.
- All log related activities (and in fact, all CC related activities such as lock/unlock, dealing with deadlocks etc.) are handled transparently by the DBMS.

Databases make these folks happy

- End users (Banks, Retailers, Scientists)
- DBMS vendors (Oracle, IBM, Microsoft)
- DB application programmers
 - Makes life easier since Dbase provides guarantees
- ✤ <u>Database administrator (DBA)</u>
 - Designs logical/physical schemas
 - Handles security and authorization
 - Data availability, crash recovery
 - Database tuning as needs evolve

Last three must understand how a DBMS works!

Fall 2010

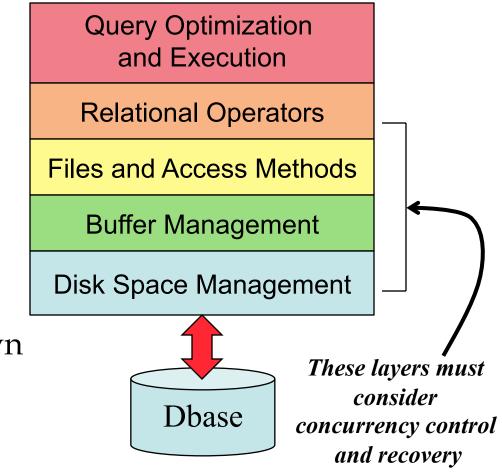






Structure of a DBMS

- A typical DBMS has a layered architecture.
- The figure does not show the concurrency control and recovery components.
- This is one of several possible architectures; each system has its own variations.





Summary



- * DBMS used to maintain, query large datasets.
- Senefits include recovery from system crashes, concurrent access, quick application development, data integrity, and security.
- * Levels of abstraction give data independence.
- * A DBMS typically has a layered architecture.
- ✤ DBAs hold responsible jobs and are well-paid! ☺
- DBMS R&D is one of the broadest, most exciting growth areas in CS.





 Modeling Data
 The Entity-Relationship (ER) model

