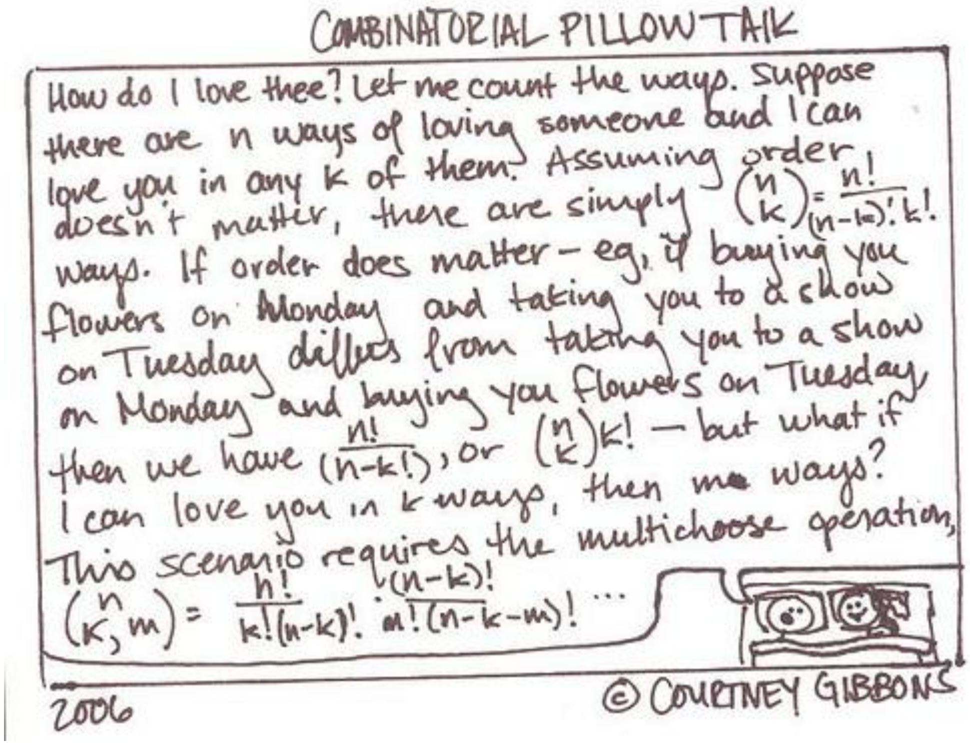
Comp 555 - BioAlgorithms - Spring 2018



- **PROBLEM SET #2 IS LATE, BUT STILL COMING**

Combinatorial Pattern Matching

A Recurring Problem

- Finding patterns within sequences

- Variants on this idea
  - Finding repeated motifs amoungst a set of strings
  - What are the most frequent k-mers
  - How many time does a specific k-mer appear

- Fundamental problem: *Pattern Matching*
  - Find all positions of a particular substring in given sequence?

Pattern Matching

- **Goal:** Find all occurrences of a pattern in a text
- **Input:** Pattern $p = p_1, p_2, \ldots p_n$ and text $t = t_1, t_2, \ldots t_m$
- **Output:** All positions $1 < i < (m – n + 1)$ such that the *n*-letter substring of t starting at i matches p

```
In [2]:  def bruteForcePatternMatching(p, t):
             locations = []
             for i in range(0, len(t)-len(p)+1):
                 if t[i:i+len(p)] == p:
                     locations.append(i)
             return locations

         print(bruteForcePatternMatching("ssi", "imissmissmississippi"))

         [11, 14]
```

Pattern Matching Performance

- Performance:
    - *m* - length of the text *t*
    - *n* - the length of the pattern *p*
    - Search Loop - executed *O(m)* times
    - Comparison - *O(n)* symbols compared
    - Total cost - *O(mn)* per pattern
- In practice, most comparisons terminate early
- Worst-case:
    - p = "AAAT"
    - t = "AAAAAAAAAAAAAAAAAAAAAAT"

We can do better!

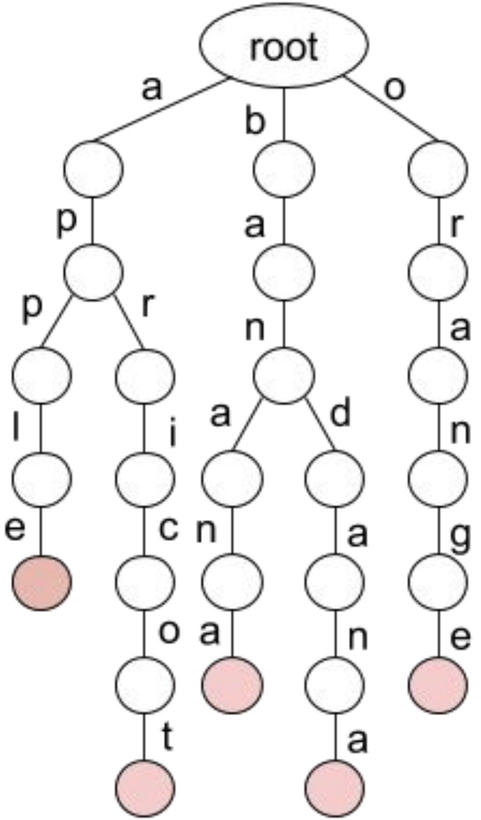
If we preprocess our pattern we can search more effciently (*O(n)*). Example:

```
         imissmissmississippi
 1.      s
 2.       s
 3.        s
 4.         SSi
 5.            s
 6.             SSi
 7.                s
 8.                 SSI             - match at 11
 9.                    SSI          - match at 14
10.                      s
11.                       s
12.                        S
```

- At steps 4 and 6 after finding the mismatch *"i" ≠ "m"* we can skip over all positions tested because we know that the suffix *"sm"* is not a prefix of our pattern *"ssi"*
- Even works for our worst-case example "AAAAT" in "AAAAAAAAAAAAAAT" by recognizing the shared prefixes ("AAA" in "AAAA").
- How about finding multiple patterns $[p_1,p_2,\ldots,p_3]$ in t

Keyword Trees

- We can preprocess the set of strings we are seeking to minimize the |number of comparisons
- **Idea:** Combine patterns that share prefixes, to *share* those comparisons
  - Stores a set of keywords in a rooted labeled tree
  - Each edge labeled with a letter from an alphabet
  - All edges leaving a given vertex have distinct labels
  - Leaf vertices are indicated
  - Every keyword stored can be spelled on a path from the root to some leaf vertex
  - Searches are performed by "threading" the target pattern through the tree
- A *Tree* is a special graph as discussed previously
  - One connected component
  - *N* nodes, *N-1* edges, No loops
  - Exactly one path from any.
- A ***Trie*** is a tree that is related to a sequence.
  - Generally, there is a 1-to-1 correspondence between either nodes or edges of the *trie* and a symbol of the sequence



Prefix *Trie* Match

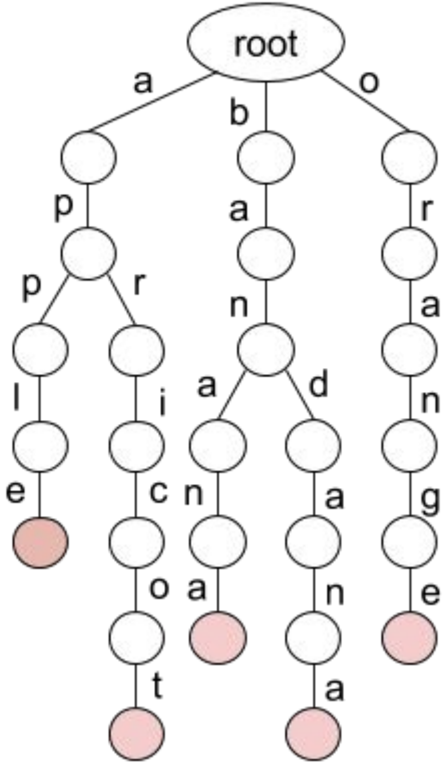
- **<u>Input:</u>** A text *t* and a trie *P* of patterns
- **<u>Output:</u>** True if *t* leads to a leaf in *P*; False otherwise

What is output for:

- *apple*
- *band*
- *april*

Performance:

- O(m) - the length of the text, *t*
- Independent of how many strings are in the Keyword Trie



Prefix *Trie* code

In [5]:

```python
def path(string, parent):
    if (len(string) > 0):
        if (string[0] in parent):
            child = parent[string[0]]
        else:
            child = {}
            parent[string[0]] = child
        path(string[1:], child)
    else:
        parent['$'] = True

class PrefixTrie:
    def __init__(self):
        """ Tree is a dictionary of the children at each node"""
        self.root = {}
    def add(self, string):
        """ Add a path from the Trie's root"""
        path(string, self.root)
    def match(self, string):
        """ Check if there is a path from the root to a '$' """
        parent = self.root
        for c in string:
            if c not in parent:
                break
            parent = parent[c]
        return '$' in parent


T = PrefixTrie()
T.add("apple")
T.add("banana")
T.add("apricot")
T.add("bandana")
T.add("orange")
print(T.root)
print([v for v in map(T.match, ['apple', 'banana', 'apricot', 'orange', 'band', 'april', 'bananapple'])])
```

```
{'a': {'p': {'p': {'l': {'e': {'$': True}}}, 'r': {'i': {'c': {'o': {'t': {'$': True}}}}}}, 'b': {'a': {'n': {'a': {'n':
{'a': {'$': True}}}, 'd': {'a': {'n': {'a': {'$': True}}}}}}}, 'o': {'r': {'a': {'n': {'g': {'e': {'$': True}}}}}}}
[True, True, True, True, False, False, True]
```

Multiple Pattern Matching

Suppose that we have a long string, *t*, like a genome, and we want to find if any of the strings in a previously constructed prefix trie, *P*, appear within it.

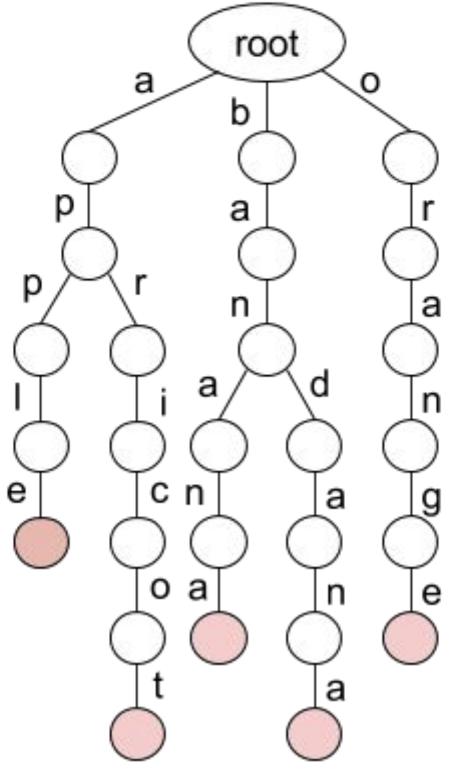
- *t* - the text to search through
- *P* - the trie of patterns to search for

```
def multiplePatternMatching(t, P):
    locations = []
    for i in xrange(0, len(t)):
        if PrefixTrieMatch(t[i:], P):
            locations.append(i)
    return locations
```

Multiple Pattern Matching Example

```
multiplePatternMatching("bananapple", P):
 0:  PrefixTrieMatching("bananapple", P) = True
 1:  PrefixTrieMatching("ananapple", P) = False
 2:  PrefixTrieMatching("nanapple", P) = False
 3:  PrefixTrieMatching("anapple", P) = False
 4:  PrefixTrieMatching("napple", P) = False
 5:  PrefixTrieMatching("apple", P) = True
 6:  PrefixTrieMatching("pple", P) = False
 7:  PrefixTrieMatching("ple", P) = False
 8:  PrefixTrieMatching("le", P) = False
 9:  PrefixTrieMatching("e", P) = False

locations = [0, 5]
```

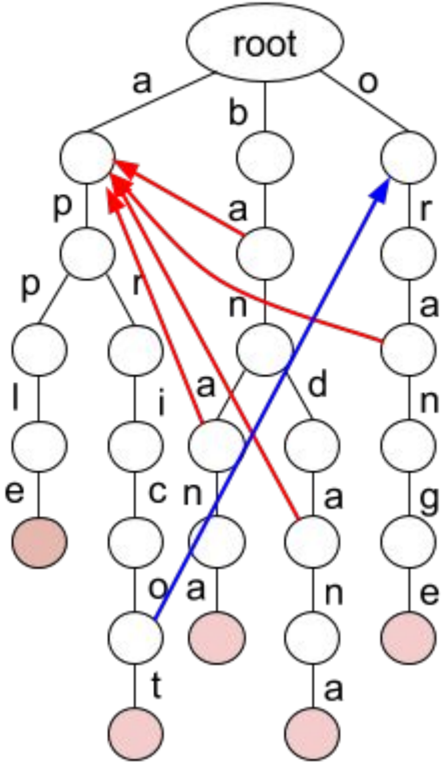


Trie Improvements

- Based on our previous speed-up
- We can add failure edges to our Trie
- *Aho-Corasick* Algorithm

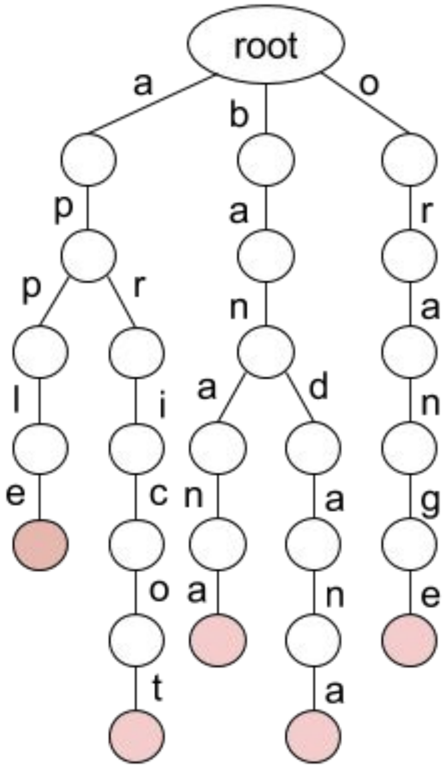
The concept of "threading" one string through another

```
bapple
bap
 apple
```



Multiple Pattern Matching Performance

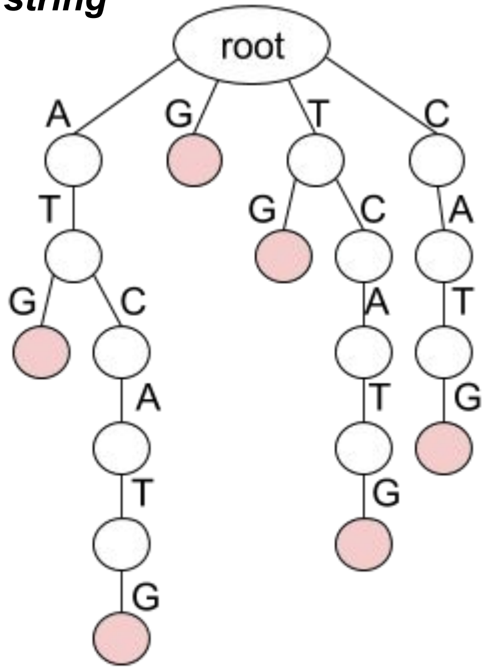
- m - len(t)
- d - max depth of P (longest pattern in P)
- O(md) to find all patterns
- Can be decreased further to O(m) using Aho-Corasick Algorithm
- Memory issues
  - Tries require a lot of memory
  - Practical implementation is challenging
  - Genomic reads - millions to billions of
- Patterns typically of length > 100



Now for a Twist

- What if our list of keywords were simply all suffixes of a ***single given string***

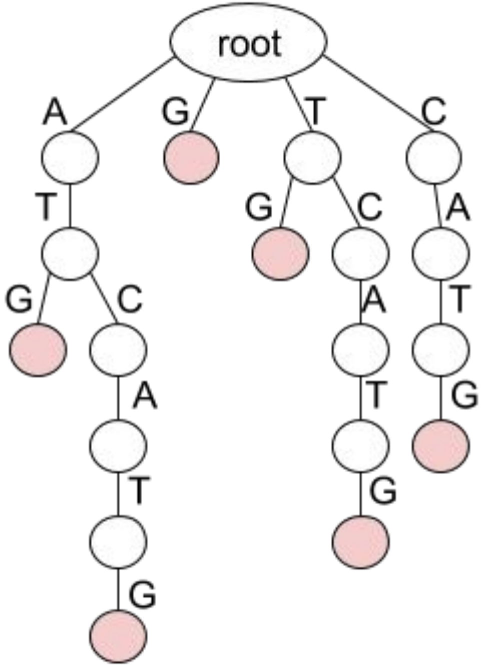
```
Example: ATCATG
          TCATG
           CATG
            ATG
             TG
              G
```



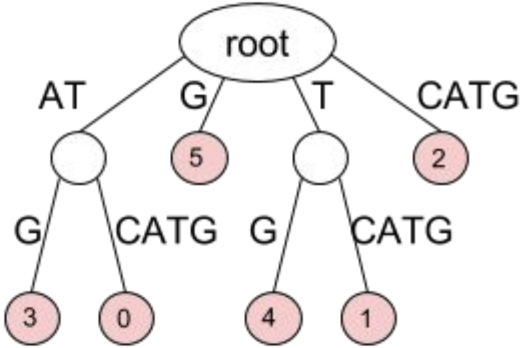
- The resulting keyword tree:
- A **Suffix Trie**

Suffix Tree

A *compressed* Suffix Trie

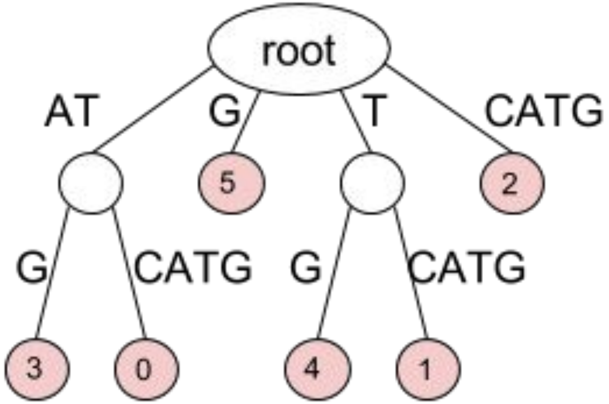


- Combine nodes with in and out degree 1
- Edges are text substrings
- All internal nodes have at least 3 edges
- All leaf nodes are labeled with an index



Uses for Suffix Trees

- Suffix trees hold all suffixes of a text, T
  - i.e., ATCATG: ATCATG, TCATG, CATG, ATG, TG, G
- Can be built in O(m) time for text of length m
- To find any pattern P in a text:
  - Build suffix tree for text, O(m), m=|T|
  - Thread the pattern through the suffix tree
  - Can find pattern in O(n) time! (n=|P|)
- O(|T|+|P|) time for "Pattern Matching Problem" (better than Naïve O(|P||T|)
- Build suffix tree and lookup pattern
- Multiple Pattern Matching in O(|T|+k|P|)

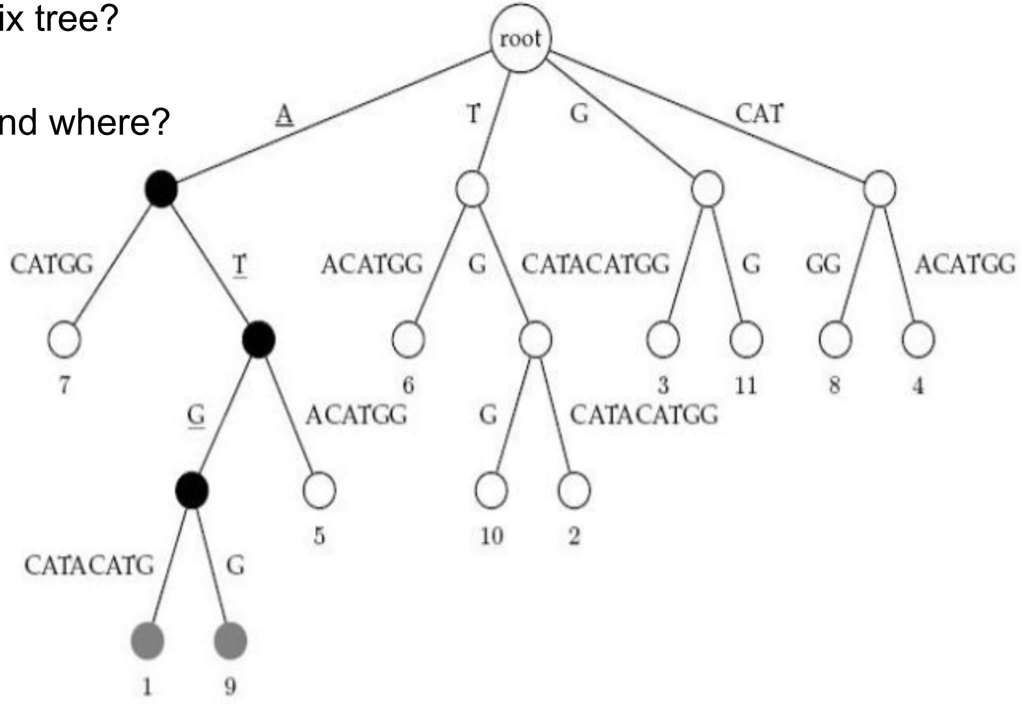


Suffix Tree Overhead

- Input: text of length m
- Computation
  - O(m) to compute a suffix tree
  - Does not require building the suffix trie first
- Memory
  - O(m) - nodes are stored as offsets and lengths
- Huge hidden constant, best implementations
- Requires about 20*m bytes
- 3 GB human genome = 60 GB RAM

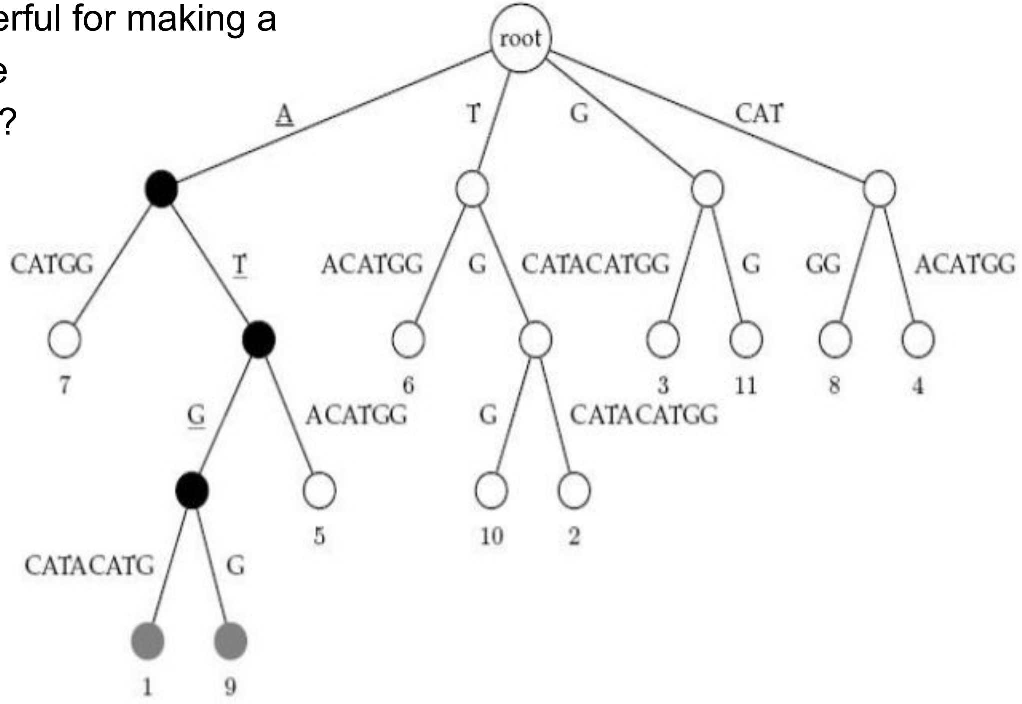
Suffix Tree Examples

- What is the string represented in the suffix tree?
- What letter occurs most frequently?
- How many times doaes "ATG" appear, and where?
- How long is the longest repeated k-mer?



Suffix Trees: Theory vs. Practice

- In theory, suffix trees are extremely powerful for making a variety of queries concerning a sequence
  - What is the shortest unique substring?
  - How many times does a given string appear in a text?
- Despite the existence of linear-time construction algorithms, and O(m) search times, suffix trees are still rarely used for genome-scale searching
- Large storage overhead



Substring Searching

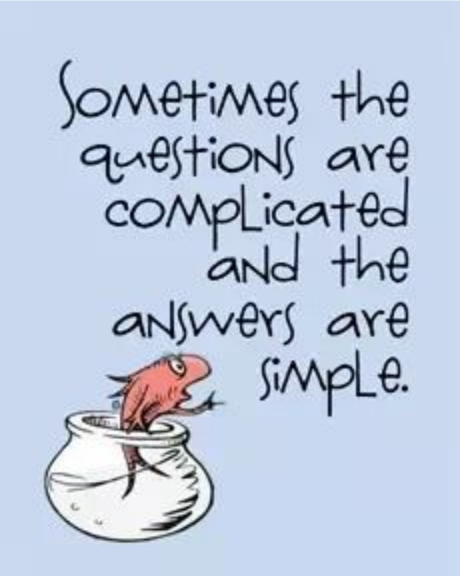
- Is there some other data structure to gain efficent access to all of the suffixes of a given string with less overhead than a suffix tree?
- Some things we know
  - Searching an unordered list of items with length *n* generally requires *O(n)* steps
  - However, if we sort our items first, then we can search using *O(log(n))* steps
  - Thus, if we plan to do frequent searches there is some advantage to performing a sort first and amortizing its cost over many searches
- For strings *suffixes* are interesting *items*. Why?

Suffixes:	**Sorted Suffixes:**
`panamabananas`	`abananas`
`anamabananas`	`amabananas`
`namabananas`	`anamabananas`
`amabananas`	`ananas`
`mabananas`	`anas`
`abananas`	`as`
`bananas`	`bananas`
`ananas`	`mabananas`
`nanas`	`namabananas`
`anas`	`nanas`
`nas`	`nas`
`as`	`panamabananas`
`s`	`s`

Questions you can ask

Is there any use for a list of sorted suffixes?

```
Sorted Suffixes:  abananas
                  amabananas
                  anamabananas
                  ananas
                  anas
                  as
                  bananas
                  mabananas
                  namabananas
                  nanas
                  nas
                  panamabananas
                  s
```



- Does the substring "nana" appear in the orginal string?
- How many times does "ana" appear in the string?
- What is the most/least frequent letter in the orginal string?
- What is the most frequent two-letter substring in the orginal string?

Properties of a sorted “suffix array”

- Size of the sorted list if the given text has a length of *m*? $O(m^2)$
- Cost of the sort? $O(m^2 \log(m))$
- Not practical for big *m*
- There are many ways to sort
    - What is an “*in place”* sort?
    - What is a “*stable”* sort?
    - What is an “*arg” sort*?

Arg Sorting

Consider the list:

```
[72,27,45,36,18,54,9,63]
```

When sorted it is simply:

```
[9,18,27,36,45,54,63,72]
```

Its “arg” sort is:

```
[6,4,1,3,2,5,7,0]
```

- The *ith* element in the arg sort is the *index* of the *ith* element from the orginal list when sorted.
- Thus, [A[i] for i in argsort(A)] == sorted[A]

Code for Arg Sorting

In [7]:

```
def argsort(input):
    return sorted(range(len(input)), key=input.__getitem__)

A = [72,27,45,36,18,54,9,63]
print(argsort(A))
print([A[i] for i in argsort(A)])

print()
B = ["TAGACAT", "AGACAT", "GACAT", "ACAT", "CAT", "AT", "T"]
print(argsort(B))
print([B[i] for i in argsort(B)])
```

```
[6, 4, 1, 3, 2, 5, 7, 0]
[9, 18, 27, 36, 45, 54, 63, 72]

[3, 1, 5, 4, 2, 6, 0]
['ACAT', 'AGACAT', 'AT', 'CAT', 'GACAT', 'T', 'TAGACAT']
```

Next Time

- We'll see how arg sorting can be used to simplify representing our sorted list of suffixes
- Suffix arrays
- Burrows-Wheeler Transforms
- Applications in sequence alignment