Perfect Phylogeny

Not in textbook
• Thus far
  – distance-based evolutionary trees
    • Additive guarantees that the tree would reproduce all pairwise distances, but not all distance matrices are additive
    • Sequences → Distances → Sequences
  – character-based evolutionary trees
    • Trees directly from sequences
    • The most general version is hard (Large parsimony)

• Infinite Sites Model
• Perfect Phylogeny
• Local vs Global Phylogenetic Trees
Character State Matrix M

- M has $n$ rows (samples)
- M has $m$ columns (characters)
- $M_{ij}$ denotes the state object $i$ has for character $j$
- Sequence Diversity Patterns (SDPs) often reoccur

same SDP
Infinite Sites Model

- Assumes mutations are rare events
- Assumes DNA sequences are large
- Multiple mutations at the same site are extremely rare
- Infinite Sites Model assumes that multiple mutations never occur at the same sequence position
- Thus, all states are "Binary" or "Biallelic"
A Different Kind of Tree

- Unrooted “Perfect Phylogeny” Tree
- Nodes correspond to sample sequences (haplotypes), both current and ancestral
- Edges correspond to actual mutations (SNPs)
- Removal of an edge creates a bipartition (each part is distinguished by a character at some position)
- SDPs can occur multiple times, and their frequency can be used as a edge weight
- Tree leaves correspond to mutations (allele variants) that are unique to a sequence, i.e. a SDP with only one minority allele instance, *private*
Unrooted Trees

• Unrooted phylogenetic trees are less specific than evolutionary trees
• The edges are undirected, thus the direction from ancestor to descendent are unknown
• All but one leaf, however, and possibly all leaves (if the root is an interior node) must be descendents
• Slightly fewer labeled unrooted trees than labeled rooted tree

\[
\frac{(2n-4)!}{2^{n-2}(n-2)!} \quad \text{vs} \quad \frac{(2n-3)!}{2^{n-2}(n-2)!}
\]

• Moreover, any node can be an “observed” sample in a phylogenetic tree whereas only leaf nodes are observed an evolutionary tree
Three different evolutionary (rooted) trees that are consistent with a common phylogenetic (unrooted) tree
Building a Phylogenetic Tree

- Assume we only have direct access to current haplotypes
- Construct a pair-wise distance matrix between haplotypes using Hamming distances
- Add smallest edge between all nodes which do not introduce a loop
- If the smallest distance is greater than 1 add d-1 “hidden” nodes between the pair so that adjacent nodes have a hamming distance of 1
- Augment the distance matrix with the new nodes and claim the introduced edges
- Repeat finding the smallest distance, and augmenting until the graph is connected
Our tree construction method will not work for any arbitrary set of character sequences; it only works for those that satisfy the assumptions of the infinite sites model.

Under the assumption of the infinite sites model all SNP pairs exhibit the property no more that 3 out of the possible 4 allele combinations occur.

Direct consequence of only one mutation per site.

Showing that all SNP pair combinations satisfy the four gamete test is a necessary and sufficient condition for there to exist a perfect phylogeny tree.

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<thead>
<tr>
<th></th>
<th>S_1</th>
<th>S_2</th>
<th>S_3</th>
<th>S_4</th>
<th>S_5</th>
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</tbody>
</table>
Questions

• Does there exist SDPs that are compatible with all others?

Private SNPs are compatible with any other SNP

• Given N distinct haplotype sequences resulting from an infinite sites model what is minimum number of SDPs?

N-1 edges are the fewest necessary to connect N haplotypes into a “linear” tree. How many singleton SNPs occur in such a tree? 2

• Given N distinct haplotype sequences resulting from an infinite sites model what is maximum number of SDPs?

2N-3 edges, the number of edges in an unrooted tree with N leaves
Exercise

Consider the following SNP panel

<table>
<thead>
<tr>
<th></th>
<th>S₁</th>
<th>S₂</th>
<th>S₃</th>
<th>S₄</th>
<th>S₅</th>
<th>S₅</th>
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<tbody>
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<tr>
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</tr>
<tr>
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<td>1</td>
<td>0</td>
<td>0</td>
</tr>
</tbody>
</table>

Satisfies the four gamete test?
Construct the tree
Is the SDP $11001^T$ possible?
Complications

• There are two issues that limit the use of Perfect Phylogeny, both are violations of our infinite-sites model assumptions
  – In addition to mutations, haplotype diversity is generated by recombination, exchange of subsequences between haplotypes
  – Mutations reoccur at the same position (Homoplasy)

• Thus, global (over the entire genome) perfect phylogenies are rare, but local perfect phylogenies are common

• How do we locate recombinations and recurrent mutations?
Non-sequence Complications

• **Evolutionary Convergence:**
  - Wings on birds and bats
  - Fins on Seals and Fish

• **Evolutionary Reversals:**
  - Fish → Lizard → Snake
  - Fish → Mammal → Manatee
  - (gain and later loss of legs)

• Such paths also violate the infinite sites model
SNP Compatibility

• How do we find local genomic regions where our assumptions are valid?
• Apply 4-gamete test
• Issues
  – Can we efficiently find all compatibility intervals
  – How many intervals? (fewest necessary to cover the entire genome)
  – Unique?
  – Common properties
Algorithms

• Left-to-right scan
• Is this solution unique?
Algorithms

- Left-to-right scan
- Is this solution unique? No.
- Right-to-Left scan

- Given that the solution is not unique, which do we choose?
- The most parsimonious
• Questions
  • Of all scans, which has the fewest intervals?
  • Is there a solution with fewer intervals?
• What is a better solution?
  – Clearly the intervals could be larger
  – What is the maximal size of the intervals?
Algorithms

• Theorem
  – Left-to-right and right-to-left scans have the same number of intervals, $k$
  – $k$ is the minimum number of intervals possible
Cores

- The interval overlaps tell us something important
  - Pair the L-R and R-L scan intervals from left to right. The overlap of these pairs are the interval cores.
    - The $i^{th}$ core essentially is the SNPs that the $i^{th}$ interval of the L-R and R-L scan agree should be included in the $i^{th}$ interval of any minimal set of intervals
  - A refinement of Parsimonious:
    - Use this to find the minimal set of maximally-sized intervals
Uber Scan

• But first, let's backup momentarily
  – The left-to-right scan found a minimal set of non-overlapping intervals
  – Can we find the set of all intervals of maximal size?
  – These were clearly not found in our left-to-right or right-to-left scans
Uber Scan

- Simple modification to the left-to-right scan algorithm
  - Instead of restarting when an incompatibility is found, only remove a portion of it
  - Specifically remove everything before (in the scanning direction) and including the closest newly introduced incompatibility
  - Open a new interval starting at the first SNP in the queue
  - Continue as before
Uber Scan

- Properties
  - Will contain more than the minimal number of intervals, $k$
  - Each interval is maximal in size (bounded on each side by an incompatibility)
  - Maintains a linear runtime
Max-\(k\) cover

• Minimal set of \(k\) maximally-sized intervals
  – Must be a subset of the Uber scan, since Uber includes all intervals of maximal size
  – Search all subsets of size \(k\)?

\[\binom{|\text{Uber}|}{k}\]

• No. Combinatorial Explosion
• Instead restructure the problem as a graph problem
Max-$k$ cover

- Minimal set of $k$ maximally-sized intervals
  - We know any minimal set must include the cores
  - Find all intervals from the Uber scan that overlap each core
  - Construct a $k$-partite graph
    - Vertices are intervals
    - Edges are weighted with the amount of overlap
  - Solve for maximal path (dynamic program)
Max-$k$ cover

- Properties
  - May not be unique
  - Theoretical runtime $O(ku)$, where $u$ is the number of intervals in Uber scan
  - In practice, we never see more than 3 intervals in any part, thus $O(k)$
Uses

- Phylogeny trees
  - Represent the data with the fewest possible trees
  - Maximal intervals provide maximal support for each tree

- Recombination
  - $k$ gives us a lower bound on the minimum number of recombinations needed to make the dataset
  - Although, not very tight
  - But it scales to large datasets
Critical SNPs

• How stable are these intervals?
  • If we remove any given SNP, will the minimal number of intervals needed, $k$, be reduced?

• Algorithm
  – Only consider the flagging SNPs of the Uber intervals
  • Intervals are bounded by incompatibilities. Unless they are removed, the interval cannot change size
Some Context

346866 of 689472 Perlegen SNPs on Chr 1, 60 Billion pairwise relationships, >7.5 GBytes
Chromosome 14
15059098-15230790

Trees based on
Perfect Phylogenies
Local to Global Trees

• Given a forest of local phylogeny trees, how do we construct a global tree?
• Generally, by combining tree metrics (Sum of distances from $i$ to $j$) across all trees and then applying either neighbor joining or UPMGA
• Evolution is more complicated than a simple tree
  – Common introgressions near species splits
  – Gene flows when branches interact